

# Appendix of "Bi-level Ensemble Method for Unsupervised Feature Selection"

## Appendix A: Proof of Theorem 1

**Theorem 1.** Denoting  $B_{ij} = \|\mathbf{P} \text{diag}(\mathbf{v}) \mathbf{x}_i - \mathbf{P} \text{diag}(\mathbf{v}) \mathbf{x}_j\|_2^2$  and  $C_{ij} = \|\mathbf{Y}_i - \mathbf{Y}_j\|_2^2$ , the closed-form solution of the subproblem w.r.t.  $\mathbf{S}$  is

$$S_{ij} = \max \left( \min \left( \frac{\sum_{k=1}^m \beta_k^2 S_{ij}^{(k)} - \frac{B_{ij} + \rho C_{ij}}{2W_{ij}^2}}{\sum_{k=1}^m \beta_k^2}, 1 \right), 0 \right). \quad (1)$$

*Proof.* When fixing other variables, by taking the definition of  $B_{ij}$  and  $C_{ij}$  into the objective function, we reformulate it as follows:

$$\begin{aligned} \min_{\mathbf{S}} \sum_{i,j=1}^n B_{ij} S_{ij} + \sum_{k=1}^m \beta_k^2 \|\mathbf{W} \odot (\mathbf{S} - \mathbf{S}^{(k)})\|_F^2 + \rho \sum_{i,j=1}^n C_{ij} S_{ij}, \\ \text{s.t. } 0 \leq S_{ij} \leq 1, \quad \mathbf{S} = \mathbf{S}^T. \end{aligned} \quad (2)$$

For simplicity, we first remove the constraint  $\mathbf{S} = \mathbf{S}^T$ , and at last we show the learned  $\mathbf{S}$  satisfies the symmetric constraint. When removing the symmetric constraint, we can decouple Eq. (2) into  $n \times n$  independent subproblems. Consider the  $(i, j)$ -th subproblem:

$$\begin{aligned} \min_{S_{ij}} \sum_{k=1}^m \beta_k^2 W_{ij}^2 (S_{ij} - S_{ij}^{(k)})^2 + (B_{ij} + \rho C_{ij}) S_{ij}, \\ \text{s.t. } 0 \leq S_{ij} \leq 1. \end{aligned} \quad (3)$$

By setting its derivation w.r.t.  $S_{ij}$  to zero, we obtain

$$S_{ij} = \frac{\sum_{k=1}^m \beta_k^2 S_{ij}^{(k)} - \frac{B_{ij} + \rho C_{ij}}{2W_{ij}^2}}{\sum_{k=1}^m \beta_k^2}. \quad (4)$$

If  $\frac{\sum_{k=1}^m \beta_k^2 S_{ij}^{(k)} - \frac{B_{ij} + \rho C_{ij}}{2W_{ij}^2}}{\sum_{k=1}^m \beta_k^2} > 1$ , Eq. (3) is a monotone decreasing function in the range

$[0, 1]$ , and thus the solution is 1. Similarly, if  $\frac{\sum_{k=1}^m \beta_k^2 S_{ij}^{(k)} - \frac{B_{ij} + \rho C_{ij}}{2W_{ij}^2}}{\sum_{k=1}^m \beta_k^2} < 0$ , the optima is

0. Therefore, the solution of  $\mathbf{S}$  is

$$S_{ij} = \max \left( \min \left( \frac{\sum_{k=1}^m \beta_k^2 S_{ij}^{(k)} - \frac{B_{ij} + \rho C_{ij}}{2W_{ij}^2}}{\sum_{k=1}^m \beta_k^2}, 1 \right), 0 \right).$$

Now, we prove that it is symmetric. we use the mathematical induction. In the first iteration, we initialize  $\mathbf{S} = \frac{1}{m} \sum_{k=1}^m \mathbf{S}^{(k)}$ , which is symmetric. Then, in the following iterations, if  $\mathbf{S}$  in the last iteration is symmetric, then according to the optimization of  $\mathbf{W}$ ,  $\mathbf{W}$  will be symmetric too. Moreover, it is easy to verify that  $\mathbf{B}$  and  $\mathbf{C}$  are always symmetric, and thus  $\mathbf{S}$  obtained by Eq. (1) satisfies the symmetric constraint.  $\square$

## Appendix B: Proof of Theorem 2

**Theorem 2.** *The following subproblem of  $\mathbf{v}$  is strictly convex quadratic programming.*

$$\begin{aligned} \min_{\mathbf{v}} \quad & \sum_{i,j=1}^n \|\mathbf{P} \text{diag}(\mathbf{v}) \mathbf{x}_i - \mathbf{P} \text{diag}(\mathbf{v}) \mathbf{x}_j\|_2^2 S_{ij} + \sum_{k=1}^m \alpha_k^2 \|\mathbf{v} - \mathbf{v}^{(k)}\|_2^2 \\ \text{s.t.} \quad & 0 \leq v_i \leq 1, \quad \sum_{i=1}^d v_i = 1. \end{aligned} \quad (5)$$

*Proof.* Obviously, Eq. (5) is a quadratic programming, and we just need to prove its convexity. The constraint is a convex set. So, we focus on the objective function. We reformulate Eq. (5) as follows:

$$\begin{aligned} \min_{\mathbf{v}} \quad & \mathbf{v}^T \left( (\mathbf{X} \mathbf{L} \mathbf{X}^T) \odot (\mathbf{P}^T \mathbf{P}) + \sum_{k=1}^m \alpha_k^2 \mathbf{I} \right) \mathbf{v} - \mathbf{v}^T \sum_{k=1}^m \alpha_k^2 \mathbf{v}^{(k)}, \\ \text{s.t.} \quad & 0 \leq v_i \leq 1, \quad \sum_{i=1}^d v_i = 1, \end{aligned} \quad (6)$$

where  $\mathbf{I}$  is an Identity matrix.

To prove its convexity, we just need to prove  $(\mathbf{X} \mathbf{L} \mathbf{X}^T) \odot (\mathbf{P}^T \mathbf{P}) + \sum_{k=1}^m \alpha_k^2 \mathbf{I}$  is a positive semi-definite (P.S.D.) matrix. Since  $\mathbf{L}$  is a Laplacian matrix, it is a P.S.D. matrix, and  $\mathbf{X} \mathbf{L} \mathbf{X}^T$  is also P.S.D. It is easy to verify that  $\mathbf{P}^T \mathbf{P}$  is also P.S.D. Then we show that the Hadamard product of two P.S.D. matrix is also P.S.D.

**Lemma 1.** *If both  $\mathbf{M}$  and  $\mathbf{N}$  are P.S.D. matrices, then  $\mathbf{M} \odot \mathbf{N}$  is also a P.S.D. matrix.*

*Proof.* If  $\mathbf{N}$  is P.S.D., we can find a matrix  $\mathbf{Q}$  such that  $\mathbf{N} = \mathbf{Q}^T \mathbf{Q}$ . Then, given any

vector  $\mathbf{x}$ , we compute  $\mathbf{x}^T(\mathbf{M} \odot \mathbf{N})\mathbf{x}$  as follows:

$$\begin{aligned}
\mathbf{x}^T(\mathbf{M} \odot \mathbf{N})\mathbf{x} &= \sum_{i,j} x_i M_{ij} N_{ij} x_j & (7) \\
&= \sum_{i,j} x_i M_{ij} \sum_k Q_{ki} Q_{kj} x_j \\
&= \sum_k (\mathbf{Q}_{k.} \odot \mathbf{x})^T \mathbf{M} (\mathbf{Q}_{k.} \odot \mathbf{x}) \\
&\geq 0
\end{aligned}$$

The last inequality is due to that  $\mathbf{M}$  is P.S.D. For any vector  $\mathbf{x}$ , we have  $\mathbf{x}^T(\mathbf{M} \odot \mathbf{N})\mathbf{x} \geq 0$ , and thus  $\mathbf{M} \odot \mathbf{N}$  is also a P.S.D. matrix.  $\square$

According to Lemma 1, we have  $(\mathbf{X}\mathbf{L}\mathbf{X}^T) \odot (\mathbf{P}^T\mathbf{P})$  is a P.S.D. matrix. Since  $\sum_{k=1}^m \alpha_k^2 > 0$ ,  $(\mathbf{X}\mathbf{L}\mathbf{X}^T) \odot (\mathbf{P}^T\mathbf{P}) + \sum_{k=1}^m \alpha_k^2 \mathbf{I}$  is a positive definite matrix. Therefore, Eq. (5) is strictly convex quadratic programming.  $\square$